

BUTTERFLY FACTORIZATION

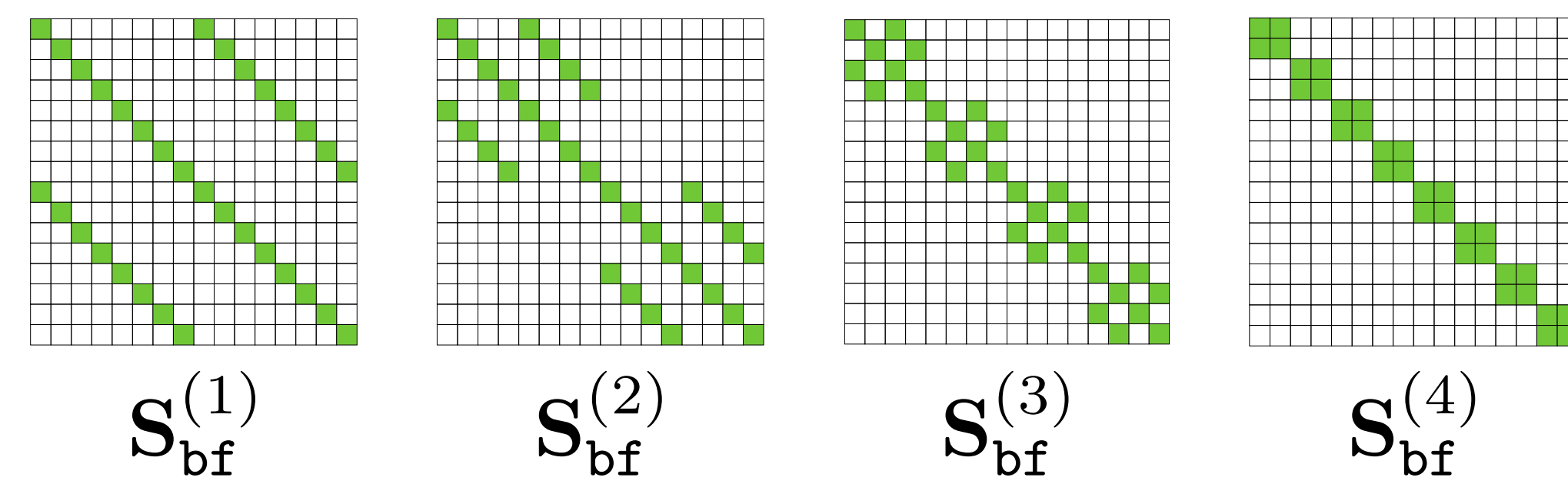
Given $\mathbf{Z} \in \mathbb{C}^{N \times N}$ with $N = 2^J$:

$$\min_{\mathbf{X}^{(1)}, \dots, \mathbf{X}^{(J)}} \|\mathbf{Z} - \mathbf{X}^{(1)} \dots \mathbf{X}^{(J)}\|_F^2$$

such that $\text{supp}(\mathbf{X}^{(\ell)}) \subseteq \mathbf{S}_{\text{bf}}^{(\ell)} := \mathbf{I}_{2^{\ell-1}} \otimes \begin{pmatrix} 1 & \\ & 1 \end{pmatrix} \otimes \mathbf{I}_{N/2^\ell}$

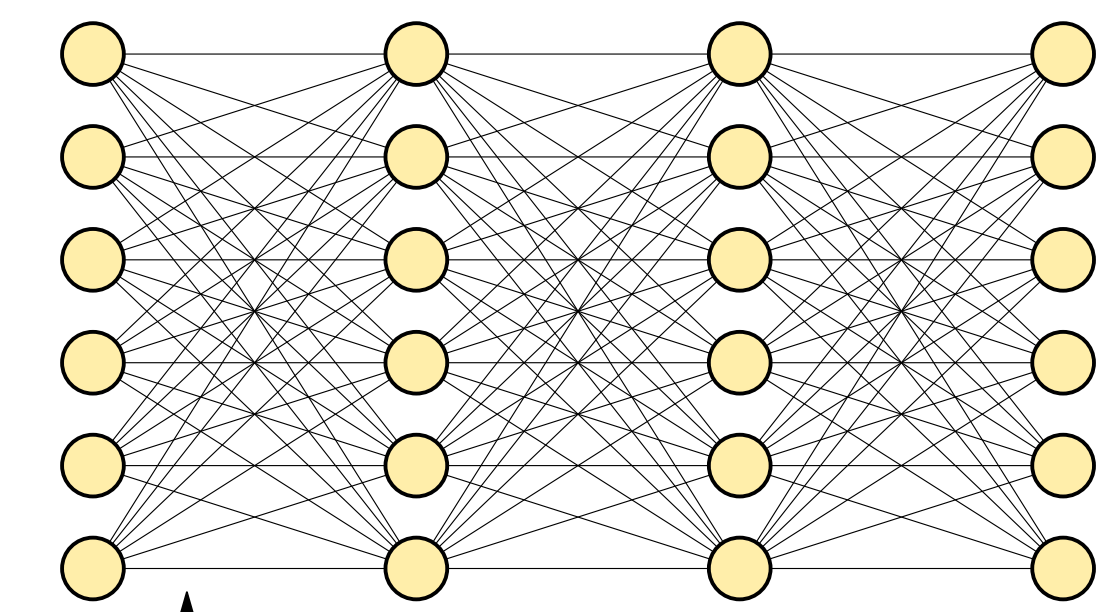
where $\text{supp}(\mathbf{M}) := \{(i, j) \mid \mathbf{M}_{i,j} \neq 0\}$.

Example with $J=4$ factors:



MOTIVATIONS

REPLACE DENSE MATRICES BY BUTTERFLY MATRICES



Less parameters
Faster training / inference time

$\mathbf{W} \approx \mathbf{X}^{(1)} \dots \mathbf{X}^{(J)}$ Fast $O(N \log N)$ matrix-vector multiplications

THE BUTTERFLY STRUCTURE IS EXPRESSIVE

$$\mathcal{B} := \{\mathbf{X}^{(1)} \dots \mathbf{X}^{(J)} \mid \text{supp}(\mathbf{X}^{(\ell)}) \subseteq \text{supp}(\mathbf{S}_{\text{bf}}^{(\ell)})\}$$

Kaleidoscope representation [2]:

$$\underbrace{\mathbf{M}^{(1,w)} \times \dots \times \mathbf{M}^{(1,w)}}_{\mathbf{M}^{(1,w)} \in \mathcal{B}} \times \underbrace{\mathbf{M}^{(2,w)*} \times \dots \times \mathbf{M}^{(2,w)*}}_{\mathbf{M}^{(2,w)*} \in \mathcal{B}^*} \in \mathcal{B}\mathcal{B}^*$$

$$\prod_{w=1}^W \mathbf{M}^{(1,w)} \mathbf{M}^{(2,w)*} \text{ captures}$$

DFT, DSC, DCT, Hadamard
Convolutions, Toeplitz
Orthogonal random features
Fastfood, ACDC

CONTRIBUTIONS

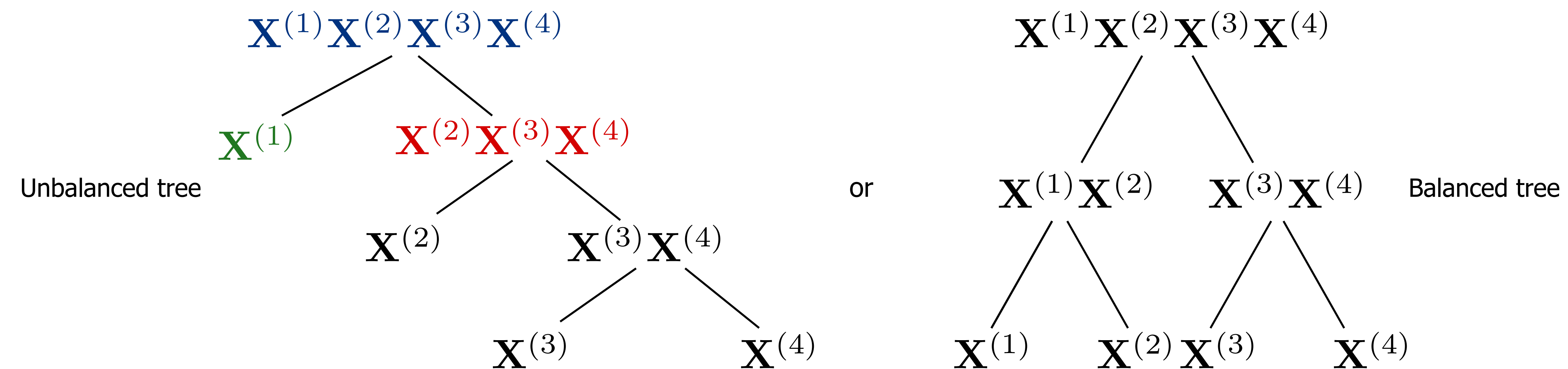
IDENTIFIABILITY

If $\mathbf{Z} = \mathbf{X}^{(1)} \times \mathbf{X}^{(2)} \times \dots \times \mathbf{X}^{(J)} = \bar{\mathbf{X}}^{(1)} \times \bar{\mathbf{X}}^{(2)} \times \dots \times \bar{\mathbf{X}}^{(J)}$
then $(\mathbf{X}^{(\ell)})_{\ell=1}^J \sim (\bar{\mathbf{X}}^{(\ell)})_{\ell=1}^J$ up to natural scaling ambiguity.

HIERARCHICAL ALGORITHM

Let $\mathbf{Z} := \mathbf{X}^{(1)} \dots \mathbf{X}^{(J)}$ with exact butterfly factorization.

Principle: perform successive two-layer fixed-support factorizations, until obtaining J factors.



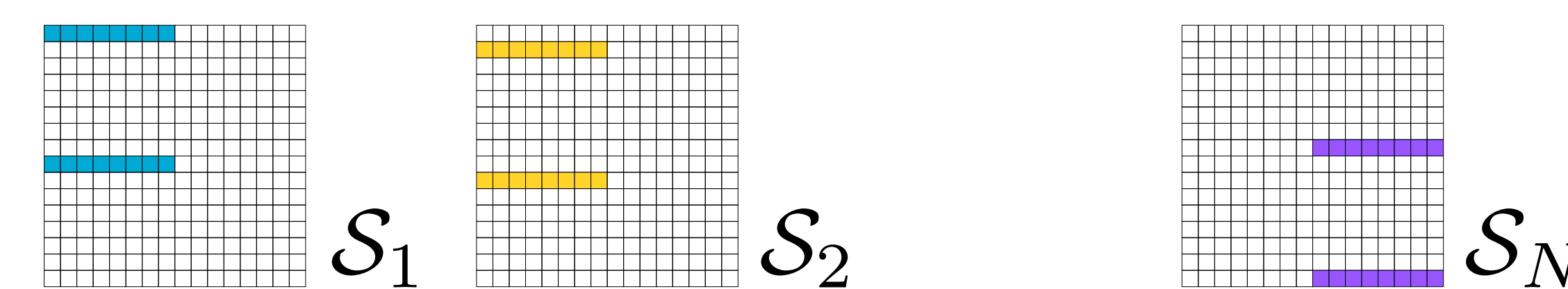
Exact recovery guarantee: the butterfly factors are recovered from \mathbf{Z} by the algorithm, in $O(N^2)$ time.

AT THE FIRST LEVEL of the unbalanced tree:

(FSMF) $\min_{\mathbf{A}, \mathbf{B}} \|\mathbf{Z} - \mathbf{A}\mathbf{B}\|_F^2$ such that $\text{supp}(\mathbf{A}) \subseteq \mathbf{S}_{\text{bf}}^{(1)}$ $\text{supp}(\mathbf{B}) \subseteq \mathbf{S}_{\text{bf}}^{(2)} \mathbf{S}_{\text{bf}}^{(3)} \mathbf{S}_{\text{bf}}^{(4)}$

Decomposition into rank-one matrices: $\mathbf{A}\mathbf{B} = \mathbf{A}_{\bullet,1}\mathbf{B}_{1,\bullet} + \mathbf{A}_{\bullet,2}\mathbf{B}_{2,\bullet} + \dots + \mathbf{A}_{\bullet,N}\mathbf{B}_{N,\bullet}$

Observation: the rank-one supports are pairwise disjoint.



THEOREM:

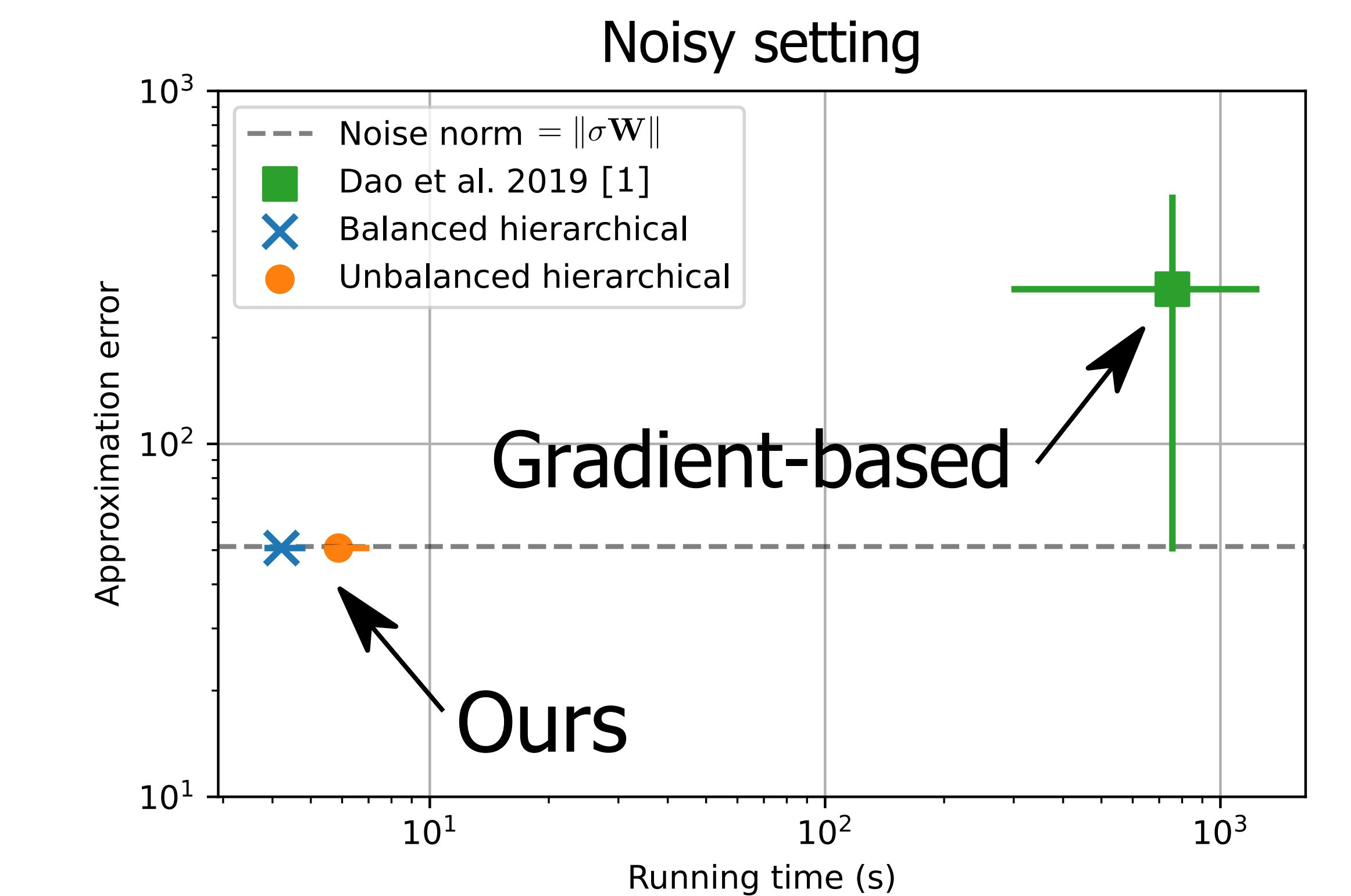
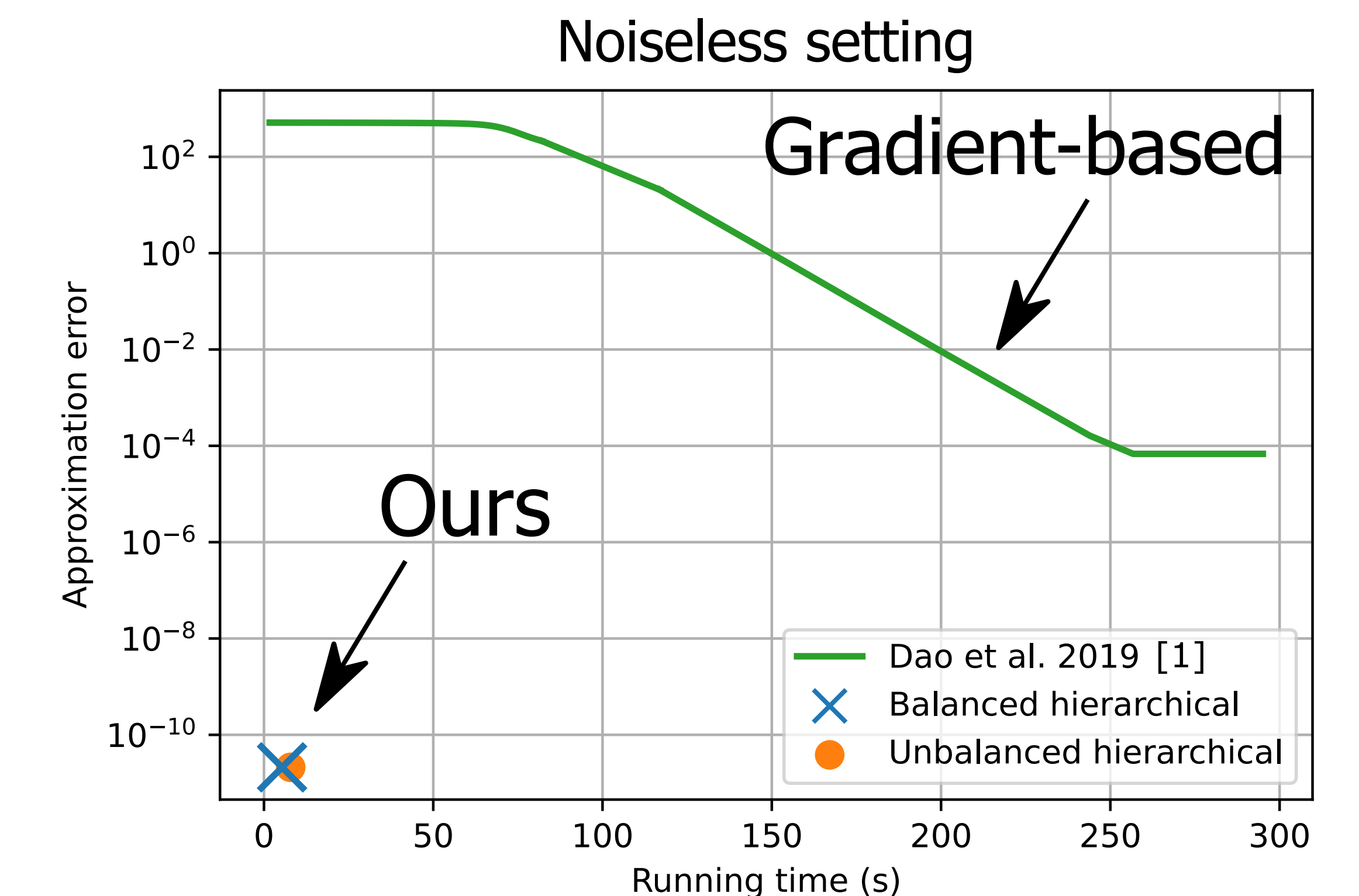
(FSMF) is polynomially solvable [3] and admits an essentially unique solution [5].

ALGORITHM to solve (FSMF) [3]:

1. Extract the submatrices $\mathbf{Z} \odot \mathcal{S}_i$ for $i = 1, \dots, N$
2. Perform best rank-one approximation for each submatrix

NUMERICAL EXPERIMENTS

Recovery of the butterfly factors of the DFT matrix:



Experiments are taken from [4].

REFERENCES

[1] T. Dao, A. Gu, M. Eichhorn, A. Rudra, and C. Ré, Learning fast algorithms for linear transforms using butterfly factorizations. In ICML, 2019.
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[3] Q.-T. Le, E. Riccietti, and R. Gribonval, Spurious Minima and NP-hardness of Sparse Matrix Factorization With Fixed Support. arXiv preprint, arXiv:2112.00386, 2021.
[4] Q.-T. Le, L. Zheng, E. Riccietti, and R. Gribonval, Fast learning of fast transforms, with guarantees. In ICASSP, 2022.
[5] L. Zheng, E. Riccietti, and R. Gribonval, Efficient Identification of Butterfly Sparse Matrix Factorizations. arXiv preprint, arXiv:2110.01230, 2022.