Efficient Identification of Butterfly Sparse Matrix Factorizations

Léon Zheng, Elisa Riccietti, Rémi Gribonval

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Factorize $\mathbf{K} := (e^{2\pi i \Phi(x,k)})_{x \in X, k \in \Omega} \approx \mathbf{X}^{(1)} \dots \mathbf{X}^{(J)}$.

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$$\min_{\mathbf{X}^{(1)},...,\mathbf{X}^{(J)}} \left\| \mathbf{Z} - \mathbf{X}^{(1)} \mathbf{X}^{(2)} ... \mathbf{X}^{(J)} \right\|_{F}, \text{ such that } \{\mathbf{X}^{(\ell)}\}_{\ell} \text{ are sparse.}$$

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When is the problem well-posed? Uniqueness of solution? Stability? \rightarrow Still an open question.

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Captures common fast transforms (Hadamard, DFT, DCT, ...)

[T. Dao et al., Kaleidoscope: An efficient, learnable representation for all structured linear maps, ICLR, 2020]

$$\mathbf{H}_{N=16} \quad \mathbf{X}^{(1)} \quad \mathbf{X}^{(2)} \quad \mathbf{X}^{(3)} \quad \mathbf{X}^{(4)}$$

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Makes the sparse matrix factorization problem well-posed [L. Zheng et al., Efficient identification of butterfly sparse matrix factorizations, SIMODS, 2023]

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• Gradient-descent method

[Le Magouarou et al., Flexible multilayer sparse approximations of matrices and applications, JSTSP, 2016] [T. Dao et al., Learning fast algorithms for linear transforms using butterfly factorizations, ICML, 2019.]

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1 We prove that the butterfly factorization is essentially unique.

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Pactors are recovered by a flexible hierarchical factorization algorithm.

The algo. is numerically faster, more accurate than gradient-descent.

Hierarchical factorization algorithm Let $\mathbf{Z} := \mathbf{X}^{(1)} \mathbf{X}^{(2)} \mathbf{X}^{(3)} \mathbf{X}^{(4)}$ such that:











How to recover the partial products?



How to recover the partial products? ightarrow use their known supports





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Two-layer fixed-support problem:

$$\min_{\mathbf{A},\mathbf{B}} \|\mathbf{Z} - \mathbf{A}\mathbf{B}\|_{F}, \text{ s.t. supp}(\mathbf{A}) \subseteq \mathbf{S}_{\mathtt{bf}}^{(1)}, \text{ supp}(\mathbf{B}) \subseteq \mathbf{S}_{\mathtt{bf}}^{(2)} \mathbf{S}_{\mathtt{bf}}^{(3)} \mathbf{S}_{\mathtt{bf}}^{(4)} \quad (\star)$$

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Constraint on the rank-one matrices

$$\sup(\mathbf{A}_{\bullet,1}\mathbf{B}_{1,\bullet}) \subseteq \mathbf{I}_{0} = \mathcal{S}_{1}$$
$$\sup(\mathbf{A}_{\bullet,2}\mathbf{B}_{2,\bullet}) \subseteq \mathbf{I}_{0} = \mathcal{S}_{2}$$
$$\vdots$$
$$\sup(\mathbf{A}_{\bullet,N}\mathbf{B}_{N,\bullet}) \subseteq \mathbf{I}_{0} = \mathcal{S}_{N}$$

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$$\operatorname{supp}(\mathbf{A}_{\bullet,1}\mathbf{B}_{1,\bullet}) \subseteq = S_1 \qquad \cdots$$
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Theorem ([Le et al. 2022; Zheng et al. 2023])

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Algorithm:

- Extract the submatrices $\mathbf{Z}_{|S_i}$, $i = 1, \dots, N$
 - Perform best rank-one approximation for each submatrix

Hierarchical factorization algorithm Let $\mathbf{Z} := \mathbf{X}^{(1)} \mathbf{X}^{(2)} \mathbf{X}^{(3)} \mathbf{X}^{(4)}$ such that: $supp(\mathbf{X}^{(1)}) \subseteq$ $supp(\mathbf{X}^{(2)}) \subseteq$ $supp(\mathbf{X}^{(2)}) \subseteq$ $supp(\mathbf{X}^{(4)}) \subseteq$

The two-layer procedure is repeated recursively.



Hierarchical factorization algorithm



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 $\{\mathbf{X}^{(\ell)}\}_{\ell=1}^4$ are recovered from **Z**, up to scaling ambiguities.

Uniqueness of butterfly factorization

Theorem ([Zheng et al. 2023])

Except for trivial degeneracies, the butterfly factorization $Z := X^{(1)} \dots X^{(J)}$ is essentially unique, up to unavoidable scaling ambiguities:

$$\begin{cases} \bar{\mathbf{X}}^{(1)} ... \bar{\mathbf{X}}^{(J)} = \mathbf{Z} \\ \forall \ell \in [J], \, \mathsf{supp}(\bar{\mathbf{X}}^{(\ell)}) \subseteq \mathbf{S}_{\mathsf{bf}}^{(\ell)} \implies (\bar{\mathbf{X}}^{(\ell)})_{\ell=1}^{J} \sim (\mathbf{X}^{(\ell)})_{\ell=1}^{J} \end{cases}$$

The unique factors are recovered from the hierarchical algorithm.

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<u>Proof</u>: at level ℓ , the intermediate matrix to factorize is

$$\mathsf{M} := \underbrace{(\mathsf{D}^{-1}\mathsf{X}^{(\ell)})}_{\mathsf{A}} \underbrace{(\mathsf{X}^{(\ell+1)} \dots \mathsf{X}^{(J)} \tilde{\mathsf{D}})}_{\mathsf{B}}.$$

The algorithm recovers (A, B) because of optimality & uniqueness in (\star) .

Flexible choice in the hierarchical order

The algorithm works for any factor-bracketing binary tree.



This extends existing work that consider only 3 trees. [Y. Liu et al., Butterfly factorization via randomized matrix-vector multiplications, SISC, 2021]

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Unbalanced tree

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Faster and more accurate in the noiseless setting

Approximate $\mathbf{Z} := \mathbf{DFT}_{512}$ by a product of J = 9 butterfly factors:



The theoretical complexity of the algorithm is $\mathcal{O}(N^2)$.

Also more robust in the noisy setting

Approximate $\mathbf{Z} := \mathbf{DFT}_{512} + \sigma \mathbf{W}$ by a product of J = 9 butterfly factors:



FAµST Toolbox: faust.inria.fr

Efficient implem. of fast transforms and algorithms for sparse matrix fact.

- Python & MATLAB wrappers (C++ core, GPU compatible)
- PYPI install: pip install pyfaust



Factorization of the Hadamard matrix with a balanced tree.

Conclusion and perspectives

- The butterfly structure captures many common fast transforms.
- **2** We proved the essential uniqueness of the butterfly factorization.
- Outputterfly factors are recovered by a flexible hierarchical algorithm.



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On going work

- Approximation error of the hierarchical algorithm
- Taking into account row and column permutations
- Efficient training of neural networks with butterfly structure

Thank you for your attention!

To know more:

- Q.-T. Le, L. Zheng, E. Riccietti, and R. Gribonval Fast learning of fast transforms, with guarantees In *ICASSP*, 2022.
- Q.-T. Le, E. Riccietti, and R. Gribonval Spurious valleys, NP-hardness, and tractability of sparse matrix factorization with fixed support In SIAM Journal on Matrix Analysis and Applications, 2022.
- L. Zheng, E. Riccietti, and R. Gribonval Efficient identification of butterfly sparse matrix factorizations In SIAM Journal on Mathematics of Data Science, 2023.